In-class Test COMS10007 Algorithms 2018/2019

12.03.2019

Reminder: $\log n$ denotes the binary logarithm, i.e., $\log n = \log_2 n$. We also write $\log^c n$ as an abbreviation for $(\log n)^c$.

Make sure to put your name on every piece of paper that you hand in!

1 *O*-notation

- 1. Let $f : \mathbb{N} \to \mathbb{N}$ be a function. Define the set $\Omega(f(n))$.
- 2. Give a formal proof of the statement:

$$10\log n \in O(\log^2 n)$$
.

- 3. For each of the following statements, indicate whether it is true of false: (no justification needed)
 - (a) $n \in O(n^2)$
 - (b) $\log n \in O(n^3)$
 - (c) $\log n \in O(\sqrt{\log n})$
 - (d) $n! \in O(2^n)$
 - (e) $2^{\sqrt{\log n}} = O(\log^2 n)$
 - (f) $f(n) \in O(g(n))$ implies $g(n) \in \Omega(f(n))$
 - (g) $f(n) \notin O(g(n))$ implies $g(n) \in O(f(n))$

2 Sorting Algorithms

Let A be an array of length n with A[i] = A[j], for every $0 \le i, j \le n - 1$.

- 1. What is the runtime of Heapsort on A? (in Θ -notation, no justification needed)
- 2. What is the runtime of Mergesort on A? (in Θ -notation, no justification needed)
- 3. What is the runtime of Insertionsort on A? (in Θ -notation, no justification needed)
- 4. What are the best-case and worst-case runtimes of Mergesort? (no justification needed)
- 5. Illustrate how the Mergesort algorithm sorts the following array (for example using a recursion tree):

 $9 \ \ 3 \ \ 2 \ \ 7 \ \ 1 \ \ 6 \ \ 11 \ \ 4$

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3 Loop-Invariant

Consider the following algorithm: (it takes two parameters, an array A of length n of positive integers, and an integer x)

Algorithm 1

Require: A is an array of n positive integers, x is an integer 1: $c \leftarrow 0$ 2: for $i \leftarrow 0, 1, ..., n-1$ do 3: if A[i] < x then 4: $c \leftarrow c+1$ 5: end if 6: end for 7: return c

1. Consider the for-loop of the algorithm. One of the following options is a correct loopinvariant:

At the beginning of iteration i (i.e., after i is updated in Line 2 and before the code in Lines 3, 4, and 5 is executed) ...

(a) ... $c = |\{j : 0 \le j < i \text{ and } A[j] < x\}|$ (b) ... $c = |\{j : 0 \le j \le i \text{ and } A[j] < x\}|$ (c) ... $c = |\{j : 0 \le j < i \text{ and } A[j] \le x\}|$ (d) ... $c = |\{j : 0 \le j \le i \text{ and } A[j] \le x\}|$

State which one is correct.

- 2. Initialization: Consider the correct invariant. Argue that at the beginning of the first iteration, i.e. when i = 0, the loop-invariant holds.
- 3. Maintenance: Consider the correct invariant. Suppose that the loop invariant holds at the beginning of iteration i. Argue that the loop-invariant then also holds at the beginning of iteration i + 1.
- 4. *Termination:* What does the algorithm compute? Argue that this follows from the correct loop invariant.