

# Advanced topics in TCS

## Exercise sheet 4.

### Minimum spanning tree, Testing $k$ -connectivity

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#### Question 1. Minimum Spanning Tree (MST)

We consider a weighted graph  $G = (V, E, w)$ , where  $w : E \rightarrow \mathbb{N}$  is an edge weight function. A *minimum spanning tree*  $F \subseteq E$  in  $G$  is a spanning tree in  $G$  of minimum weight, i.e., the sum of its edge weights is as small as possible.

We consider the streaming edge-arrival model where the edges arrive together with their weights. More specifically, the input stream consists of a sequence of tuples  $(e_i, w(e_i))_i$ , where  $w(e_i)$  is the weight of edge  $e_i$ .

1. Give a 1-pass semi-streaming algorithm for computing an MST.

*Hint:* Adapt the spanning tree algorithm from the lecture.

2. Let  $E_i$  be the first  $i$  edges in the stream,  $G_i = (V, E_i, w|_{E_i})$  (where  $w|_{E_i}$  denotes the weight function  $w$  restricted to the domain  $E_i$ ), and let  $F_i$  denote the collection of edges stored by the algorithm given in the previous exercise after iteration  $i$ . Prove by induction that  $F_i$  is a MST in  $G_i$ .

The following property may be useful:

**Lemma 1.** *Let  $T \subseteq E$  be a spanning tree in a weighted graph  $G = (V, E, w)$ . Then, if  $T$  is not a minimum spanning tree, then there exists an edge  $e \in E \setminus T$  such that  $w(e) < w(f)$ , for at least one edge  $f$  different to  $e$  in the unique cycle in  $T \cup \{e\}$ .*

#### Question 2. Deciding $k$ -Connectivity

We say that a graph  $G$  is  *$k$ -connected* if we need to remove at least  $k$  edges from  $G$  in order to disconnect  $G$ .

Consider the following algorithm for deciding  $k$ -connectivity of a graph:

1.  $F_1, F_2, \dots, F_k \leftarrow \emptyset$
2. For each edge  $e$  in the stream: If there is an  $i \in \{1, \dots, k\}$  such that  $F_i \cup \{e\}$  has no cycle then add  $e$  to  $F_i$  (if there are multiple such  $i$  then pick only one, ties can be broken arbitrarily)
3. Post-processing: Let  $F = \bigcup_{i=1}^k F_i$   
If  $(V, F)$  is  $k$ -connected then **return** “ $G$  is  $k$ -connected”, otherwise **return** “ $G$  is not  $k$ -connected”

Algorithm 1.

1. How much space does Algorithm 1 use (as a function of  $n$  and  $k$ )?
2. Prove that the algorithm is correct.