

```
% search(Agenda,Goal) <- Goal is a goal node, and a
%                               descendant of one of the nodes
%                               on the Agenda
search(Agenda,Goal):-
  next(Agenda,Goal,Rest),
  goal(Goal).
search(Agenda,Goal):-
  next(Agenda,Current,Rest),
  children(Current,Children),
  add(Children,Rest,NewAgenda),
  search(NewAgenda,Goal).
```

```
search_df([Goal|Rest],Goal):-
    goal(Goal).
search_df([Current|Rest],Goal):-
    children(Current,Children),
    append(Children,Rest,NewAgenda),
    search_df(NewAgenda,Goal).

search_bf([Goal|Rest],Goal):-
    goal(Goal).
search_bf([Current|Rest],Goal):-
    children(Current,Children),
    append(Rest,Children,NewAgenda),
    search_bf(NewAgenda,Goal).

children(Node,Children):-
    findall(C,arc(Node,C),Children).
```

Depth-first vs. breadth-first search

 Breadth-first search

- ✓ agenda = queue (first-in first-out)
- ✓ complete: guaranteed to find all solutions
- ✓ first solution founds along shortest path
- ✓ requires $O(B^n)$ memory

 Depth-first search

- ✓ agenda = stack (last-in first-out)
- ✓ incomplete: may get trapped in infinite branch
- ✓ no shortest-path property
- ✓ requires $O(B \times n)$ memory



Depth-first vs. breadth-first search

```
% depth-first search with loop detection
search_df_loop([Goal|Rest], Visited, Goal):-
    goal(Goal).
search_df_loop([Current|Rest], Visited, Goal):-
    children(Current, Children),
    add_df(Children, Rest, Visited, NewAgenda),
    search_df_loop(NewAgenda, [Current|Visited], Goal).

add_df([], Agenda, Visited, Agenda).
add_df([Child|Rest], OldAgenda, Visited, [Child|NewAgenda]):-
    not element(Child, OldAgenda),
    not element(Child, Visited),
    add_df(Rest, OldAgenda, Visited, NewAgenda).
add_df([Child|Rest], OldAgenda, Visited, NewAgenda):-
    element(Child, OldAgenda),
    add_df(Rest, OldAgenda, Visited, NewAgenda).
add_df([Child|Rest], OldAgenda, Visited, NewAgenda):-
    element(Child, Visited),
    add_df(Rest, OldAgenda, Visited, NewAgenda).
```

```
% depth-first search by means of backtracking
search_bt(Goal,Goal):-
    goal(Goal).
search_bt(Current,Goal):-
    arc(Current,Child),
    search_bt(Child,Goal).

% backtracking depth-first search with depth bound
search_d(D,Goal,Goal):-
    goal(Goal).
search_d(D,Current,Goal):-
    D>0, D1 is D-1,
    arc(Current,Child),
    search_d(D1,Child,Goal).
```

Backtracking search

```
search_id(First,Goal):-  
    search_id(1,First,Goal).           % start with depth 1  
  
search_id(D,Current,Goal):-  
    search_d(D,Current,Goal).  
search_id(D,Current,Goal):-  
    D1 is D+1,                         % increase depth  
    search_id(D1,Current,Goal).
```

☞ combines advantages of
breadth-first search (complete, shortest path)
with those of depth-first search (memory-efficient)

```
prove_df_a(Goal):-
  prove_df_a([Goal]).

prove_df_a([true|Agenda]).
prove_df_a([(A,B)|Agenda]):-!,
  findall(D,(clause(A,C),conj_append(C,B,D)),Children),
  append(Children,Agenda,NewAgenda),
  prove_df_a(NewAgenda).
prove_df_a([A|Agenda]):-
  findall(B,(clause(A,B),Children),
  append(Children,Agenda,NewAgenda),
  prove_df_a(NewAgenda).

prove(true):-!.
prove((A,B)):-!,
  clause(A,C),
  conj_append(C,B,D),
  prove(D).
prove(A):-
  clause(A,B),
  prove(B).
```

```
refute((false:-true)).
refute((A,C)):-
    cl(Cl),
    resolve(A,Cl,R),
    refute(R).

% refute_bf(Clause) <- Clause is refuted by clauses
%                       defined by cl/1
%                       (breadth-first search strategy)
refute_bf_a(Clause):-
    refute_bf_a([a(Clause,Clause)],Clause).

refute_bf_a([a((false:-true),Clause)|Rest],Clause).
refute_bf_a([a(A,C)|Rest],Clause):-
    findall(a(R,C),(cl(Cl),resolve(A,Cl,R)),Children),
    append(Rest,Children,NewAgenda),% breadth-first
    refute_bf_a(NewAgenda,Clause).
```

Refutation prover for clausal logic


```
% model(M) <- M is a model of the clauses defined by cl/1
model(M):-
  model([],M).

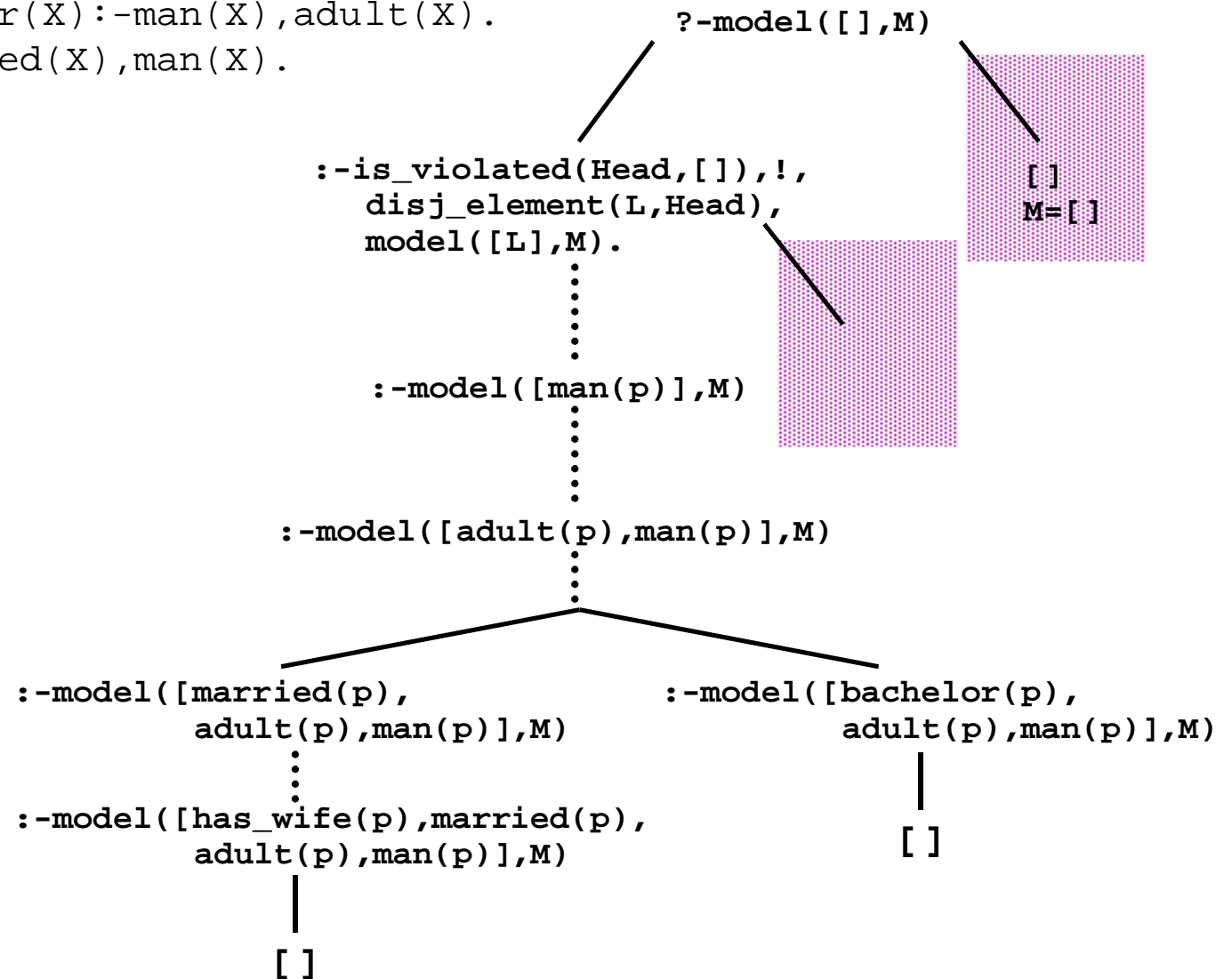
model(M0,M):-
  is_violated(Head,M0),!,           % instance of violated clause
  disj_element(L,Head),           % L: ground literal from head
  model([L|M0],M).                % add L to the model
model(M,M).                        % no more violated clauses

is_violated(H,M):-
  cl((H:-B)),                      % grounds the variables
  satisfied_body(B,M),
  not satisfied_head(H,M).
```

```

married(X);bachelor(X):-man(X),adult(X).
has_wife(X):-married(X),man(X).
man(paul).
adult(paul).

```



Forward chaining: example

```
% model_d(D,M) <- M is a submodel of the clauses
%               defined by cl/1
model_d(D,M):-
    model_d(D,[],M).

model_d(0,M,M).
model_d(D,M0,M):-
    D>0,D1 is D-1,
    findall(H,is_violated(H,M0),Heads),
    satisfy_clauses(Heads,M0,M1),
    model_d(D1,M1,M).

satisfy_clauses([],M,M).
satisfy_clauses([H|Hs],M0,M):-
    disj_element(L,H),
    satisfy_clauses(Hs,[L|M0],M).
```

Forward chaining with depth-bound